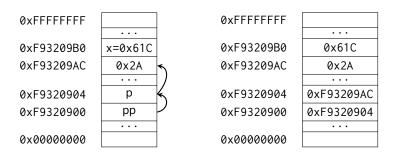
# 1 C

C is syntactically similar to Java, but there are a few key differences:

- 1. C is function-oriented, not object-oriented; there are no objects.
- 2. C does not automatically handle memory for you.
  - Stack memory, or *things that are not manually allocated*: data is garbage immediately after the *function in which it was defined* returns.
  - Heap memory, or *things allocated with* malloc, calloc, *or* realloc: data is freed only when the programmer explicitly frees it!
  - There are two other sections of memory that we learn about in this course, *static* and *code*, but we'll get to those later.
  - In any case, allocated memory always holds garbage until it is initialized!
- 3. C uses pointers explicitly. If p is a pointer, then \*p tells us to use the value that p points to, rather than the value of p, and &x gives the address of x rather than the value of x.

On the left is the memory represented as a box-and-pointer diagram.

On the right, we see how the memory is really represented in the computer.



Let's assume that int \* p is located at 0xF9320904 and int x is located at 0xF93209B0. As we can observe:

- \*p evaluates to 0x2A (42<sub>10</sub>).
- p evaluates to 0xF93209AC.
- x evaluates to 0x61C.
- &x evaluates to 0xF93209B0.

Let's say we have an int \*\*pp that is located at 0xF9320900.

#### 2 C Basics

[1.1] What does pp evaluate to? How about \*pp? What about \*\*pp?

pp evaluates to 0xF9320904. \*pp evaluates to 0xF93209AC. \*\*pp evaluates to 0x2A.

1.2 The following functions are syntactically-correct C, but written in an incomprehensible style. Describe the behavior of each function in plain English.

(a) Recall that the ternary operator evaluates the condition before the ? and returns the value before the colon (:) if true, or the value after it if false.

```
int foo(int *arr, size_t n) {
    return n ? arr[0] + foo(arr + 1, n - 1) : 0;
}
```

Returns the sum of the first N elements in arr.

(b) Recall that the negation operator, !, returns 0 if the value is non-zero, and 1 if the value is 0. The ~ operator performs a *bitwise not* (NOT) operation.

```
int bar(int *arr, size_t n) {
    int sum = 0, i;
    for (i = n; i > 0; i--)
        sum += !arr[i - 1];
    return ~sum + 1;
    }
```

Returns -1 times the number of zeroes in the first N elements of arr.

(c) Recall that  $\widehat{}$  is the  $bitwise\ exclusive-or\ (XOR)$  operator.

```
void baz(int x, int y) {
    x = x ^ y;
    y = x ^ y;
    x = x ^ y;
    x = x ^ y;
    y;
```

Ultimately does not change the value of either x or y.

## 2 Programming with Pointers

2.1 Implement the following functions so that they work as described.

(a) Swap the value of two **int**s. *Remain swapped after returning from this function*.

```
void swap(int *x, int *y) {
    int temp = *x;
    *x = *y;
    *y = temp;
  }
```

(b) Return the number of bytes in a string. Do not use strlen.

```
int mystrlen(char* str) {
    int count = 0;
    while (*str++) {
        count++;
    }
    }
    return count;
    }
```



The following functions may contain logic or syntax errors. Find and correct them.

(a) Returns the sum of all the elements in summands.

It is necessary to pass a size alongside the pointer.

```
int sum(int* summands, size_t n) {
    int sum = 0;
    for (int i = 0; i < n; i++)
        sum += *(summands + i);
    return sum;
    }
</pre>
```

(b) Increments all of the letters in the string which is stored at the front of an array of arbitrary length, n >= strlen(string). Does not modify any other parts of the array's memory.

The ends of strings are denoted by the null terminator rather than n. Simply having space for n characters in the array does not mean the string stored inside is also of length n.

```
void increment(char* string) {
    for (i = 0; string[i] != 0; i++)
    string[i]++; // or (*(string + i))++;
}
```

Another common bug to watch out for is the corner case that occurs when incrementing the character with the value 0xFF. Adding 1 to 0xFF will overflow back to 0, producing a null terminator and unintentionally shortening the string.

(c) Copies the string src to dst.

```
void copy(char* src, char* dst) {
    while (*dst++ = *src++);
    }
}
```

No errors.

(d) Overwrites an input string src with "61C is awesome!" if there's room. Does nothing if there is not. Assume that length correctly represents the length of src.

```
void cs61c(char* src, size_t length) {
```

```
2 char *srcptr, replaceptr;
```

### 4 C Basics

```
char replacement[16] = "61C is awesome!";
3
        srcptr = src;
4
        replaceptr = replacement;
5
        if (length >= 16) {
6
             for (int i = 0; i < 16; i++)</pre>
7
                 *srcptr++ = *replaceptr++;
8
        }
9
10
   }
```

char \*srcptr, replaceptr initializes a char pointer, and a char—not two
char pointers.

The correct initialization should be, **char** \*srcptr, \*replaceptr.

### 3 Memory Management

3.1

For each part, choose one or more of the following memory segments where the data could be located: **code**, **static**, **heap**, **stack**.

(a) Static variables

Static

(b) Local variables

Stack

(c) Global variables

Static

(d) Constants

Code, static, or stack

Constants can be compiled directly into the code. x = x + 1 can compile with the number 1 stored directly in the machine instruction in the code. That instruction will always increment the value of the variable x by 1, so it can be stored directly in the machine instruction without reference to other memory. This can also occur with pre-processor macros.

Constants can also be found in the stack or static storage depending on if it's declared in a function or not.

```
1 const int x = 1;
2
3 int sum(int* arr) {
```

```
4 int total = 0;
5 ...
```

6 }

In this example, x is a variable whose value will be stored in the static storage, while total is a local variable whose value will be stored on the stack. Variables declared **const** are not allowed to change, but the usage of **const** can get more tricky when combined with pointers.

(e) Machine Instructions

Code

 $(f) \ {\rm Result} \ of \ {\tt malloc}$ 

Heap

(g) String Literals

Static or stack.

When declared in a function, string literals can be stored in different places. **char**\* s = "string" is stored in the static memory segment while **char**[7] s = "string" will be stored in the stack.

- 3.2 Write the code necessary to allocate memory on the heap in the following scenarios
  - (a) An array arr of k integers

arr = (int \*) malloc(sizeof(int) \* k);

(b) A string str containing p characters

str = (char \*) malloc(sizeof(char) \* (p + 1)); Don't forget the null terminator!

(c) An  $n \times m$  matrix mat of integers initialized to zero.

```
mat = (int *) calloc(n * m, sizeof(int));
```

Alternative solution. This might be needed if you wanted to efficiently permute the rows of the matrix.

```
1 mat = (int **) calloc(n, sizeof(int *));
2 for (int i = 0; i < n; i++)
3 mat[i] = (int *) calloc(m, sizeof(int));
```

Suppose we've defined a linked list **struct** as follows. Assume **\*lst** points to the first element of the list, or is NULL if the list is empty.

```
struct ll_node {
    int first;
    struct ll_node* rest;
}
```

3.3 Implement prepend, which adds one new value to the front of the linked list. Hint: why use ll\_node \*\* *lst* instead of ll\_node\**lst*?

```
void prepend(struct ll_node** lst, int value) {
    struct ll_node* item = (struct ll_node*) malloc(sizeof(struct ll_node));
    item->first = value;
    item->rest = *lst;
        *lst = item;
    }
```

[3.4] Implement free\_11, which frees all the memory consumed by the linked list.

```
void free_ll(struct ll_node** lst) {
    if (*lst) {
        free_ll(&((*lst)->rest));
        free(*lst);
        }
        *lst = NULL; // Make writes to **lst fail instead of writing to unusable memory.
    }
```